# Comparison of Minisatellites 

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## Outline

1. Tandem repeats evolution, and minisatellite maps.
2. Evolutionary model.
3. Arches.
4. Alignment algorithm.
5. Applications.

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## Minisatellites (ms): what are they?

- Class of tandem repeat: Satellites, Minisatellites, Microsatellites.
- Example of a tandem repetition:
... cggcgat cggcgac cggagat cggcgat cggcgat cggagat cgacgat ...
- A minisatellite unit measures between $7-100 \mathrm{bp}$, and their overall length is greater than 0.5 kb .
- ms undergo amplifications and contractions $\Rightarrow$ variation in their number of unit.
- ms also undergo point-mutations and homogenization.


## Interests in Minisatellites

- Implication in diseases: insulin-dependent diabetes, several cancers, epilepsy, and others [Buard, Jeffreys 97].
- Phylogeny and population studies:
- Micro-evolution in small range of evolutionary time.
- Migration of population: out-of-Africa hypothesis [Armour et al 96].
- Evolution of the Y chromosome [Jobling, Tyler 00].
- Individual or species identification: forensic studies, genetic markers, identification of bacteria [Jeffreys 93]
- Theoretical interest: Why such sequences exist? How do they evolve?


## Minisatellites maps

- Difficulty to sequence and assemble minisatellites.
- There is a specific method to obtain the sequence of variants of the unit, Minisatellite Variant Repeat PCR [Jeffreys et al. 91].
MVR-PCR yields a ms map: a sequence of symbols, each representing a different variant of the repeated unit.
- Example of a minisatellite map:
- $\mathrm{s}=\mathrm{c} g \mathrm{gc}$ gat cggcgac cggagat cggcgat cggcgat cggagat cgacgat
- New alphabet: $\mathrm{A}=\operatorname{cggcgat} \mathrm{B}=\operatorname{cggcgac} \mathrm{C}=$ cggagat $\mathrm{D}=$ cgacgat
- Corresponding map: A B C A A C D


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## Example of evolution of ms

Ancestor map: AAAAA
Today's variants: A,B,C,D,E

Trace of the evolution in individual 1:

|  | A | A | A | A | A |  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Mutation | A | A | A | B | A |  |  |  |  |  |  |
| 5 Amplifications | A | A | A | B | B | B | B | B | B | A |  |
| Mutation | A | A | A | B | B | C | B | B | B | A |  |
| Mutation | A | A | A | B | B | C | B | D | B | A |  |
| Amplification | A | A | A | B | B | C | B | D | D | B | A |

## Example of evolution of ms (2)

Trace of the evolution in individual 2:

|  | A | A | A | A | A |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Mutation | A | E | A | A | A |  |
| Amplification | A | E | A | A | A | A |
| Contraction | A | E | A | A | A |  |

We would like to compute an optimal alignment:

| I1: | A | A | A | B | B | C | B | D | D | B | A |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mid$ | C | $\mid$ | $[$ | 1 |  | $/$ |  | $/$ | $/$ | $\mid$ |
| I2: | A | E | A | A | - | - | - | - | - | - | A |

## Single Step Evolutionary Model

- Unitary operations
- Amplification(A)/Contraction(C) duplicates/removes a variant which is next to an identical variant.
- Mutation(M) mutates one variant in another.
- Insertion(I)/Deletion(D) inserts/removes a variant, but without constraint.
- Non-unitary operations:
- Generation(G)/Compression(K) generates/compresses an arch from/to its single ancestor variant.
- A cost is associated to each operation. Alignment cost is the sum of its operations costs.


## Operations costs

- The model is symmetric: $\mathrm{I}=\mathrm{D}$ and $\mathrm{A}=\mathrm{C}$.
- Observed much higher relative frequency of amplifications and contractions compared to other events $\Rightarrow \mathrm{A}, \mathrm{C}<\mathrm{M}, \mathrm{D}, \mathrm{I}$.
- A deletion can also be obtained by a mutation plus a contraction. We have either:

1. Hypothesis 1 (H1): $\mathrm{D}>\mathrm{M}+\mathrm{C}$ and $\mathrm{I}>\mathrm{A}+\mathrm{M}$
2. Hypothesis 2 (H2): $\mathrm{D} \leq \mathrm{M}+\mathrm{C}$ and $\mathrm{I} \leq \mathrm{A}+\mathrm{M}$

- Theorem: Under H1 or H2, alignment cost is a metric distance.


## Problem defi nition

Problem: Let s and r be two ms maps of length n and m , fi nd an optimal global alignment between s and r under the single step evolutionary model.

## Related works

- Alignment with duplication [Benson 97]:
- Alignment of two sequences that may contain several tandem repeats; operations are substitutions, indels and duplications.
- Differences with our approach: duplication relates the two sequences while for us a duplication is internal to one map; repeated units are unknown.
- Duplication history reconstruction [Benson et Dong 99, Elemento al 01, Tang al 01, Elemento 02, Jaitly 02].


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## Arches

Let s be a map.

Def: an arch of $s$ is a subsequence of $s$ whose first and last variants are identical.

Def: an arch is simple if its internal variants occur only once and complex otherwise.

Ex: s=acacdefda


Def: two arches are compatible if they do not cross or share a same first foot or a same last foot.

## Generate/Compress an Arch

- Generation and Compression are symmetrical.
- Generate a simple arch: from A to ABCA

| First method |  |  |  | Second method |  |  |  |  |
| :--- | :---: | :--- | :--- | :--- | :--- | :--- | :--- | :---: |
|  |  |  |  |  |  |  |  |  |
|  | A |  |  | A |  |  |  |  |
| Amplification | A | A |  |  | Amplification | A | A |  |
|  |  |  |  |  |  |  |  |  |
| Mutation | A | B |  |  |  |  |  |  |
| Amplification | A | B | B |  | Amplification | A | A |  |
| A A |  |  |  |  |  |  |  |  |
| Mutation | A | B | C |  | Mutation | A | B |  |
| A | A | A |  |  |  |  |  |  |
| Amplification | A | B | C | C | Mutation | A | B |  |
| C | A | A |  |  |  |  |  |  |
| Mutation | A | B | C | A |  |  |  |  |
|  |  |  |  |  |  |  |  |  |

Identical variants at the extremities allow to avoid one mutation.

- Non-commutativity of operations.
- Theorem: the optimal generation/compression of an arch uses the largest subset of compatible inner arches.


## Overlap graphs

- Map of size n , arches are subintervals in [1,n].
- Compatibility relation defines an overlap graph.
- Theorem: Finding the largest subset of compatible inner arches is equivalent to finding the max stable in an overlap graph.
- Complexity of this procedure is $\mathrm{O}\left(\mathrm{V}^{2}\right)$ with V the number of vertices of the graph [Apostolico et al. 92].
- Maximal number of arches is $\Theta\left(\mathrm{n}^{2}\right)$.


## Example



How to convert the problem of finding the max subset of compatible arches in the problem of finding the max stable set in an overlap graph

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## Dynamic Programming Dependencies



$$
\mathscr{A}[i, j]=\min \left\{\begin{array}{lr}
\mathscr{A}[\mathrm{i}-1, \mathrm{j}-1]+\mathrm{M}(\mathrm{~s}[\mathrm{i}], \mathrm{r}[\mathrm{j}]) & \text { Mute } \\
\mathscr{A}[\mathrm{i}-1, \mathrm{j}]+\mathrm{C} & \text { Contract } \\
\text { iff } \mathrm{s}[\mathrm{i}-1]=\mathrm{s}[\mathrm{i}] \text { or } \mathrm{s}[\mathrm{i}]=\mathrm{r}[\mathrm{j}] & \\
\mathscr{A}[\mathrm{i}-1, \mathrm{j}]+\mathrm{S} & \text { Delete } \\
\mathscr{A}[\mathrm{i}, \mathrm{j}-1]+\mathrm{A} \\
\text { iff } \mathrm{r}[\mathrm{j}-1]=\mathrm{r}[\mathrm{j}] \text { or } \mathrm{s}[\mathrm{i}]=\mathrm{r}[\mathrm{j}] & \\
\mathscr{A}[\mathrm{i}, \mathrm{j}-1]+\mathrm{I} & \text { Amplify } \\
\mathscr{A}[1, \mathrm{j}]+\mathrm{K}(\mathrm{~s}[1, \mathrm{i}]) \text { Compress an arch } \\
\forall \mathrm{l} \in[1, \mathrm{i}-2] \text { such that } \mathrm{s}[\mathrm{l}]=\mathrm{s}[\mathrm{i}] \\
\mathscr{A}\left[\mathrm{i}, \mathrm{l}^{\prime}\right]+\mathrm{G}\left(\mathrm{r}\left[\mathrm{l}^{\prime}, j\right]\right) \text { Generate an arch } \\
\forall \mathrm{l}^{\prime} \in[1, \mathrm{j}-2] \text { such that } \mathrm{r}\left[\mathrm{l}^{\prime}\right]=\mathrm{r}[\mathrm{j}]
\end{array}\right.
$$

## Complexity

Theorem: If $\mathrm{p}=\max (\mathrm{n}, \mathrm{m})$, overall complexity is $\mathrm{O}\left(\mathrm{p}^{4}\right)$ time.

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## Minisatellite MSY1

- MSY1 is a ms on the Y chromosome. Unit length $25 \mathrm{bp}, 5$ variants differ by at most 4 residues. Evolution simplified: no exchange between alleles [Jobling et al 98].
- Maps for 690 men taken in $\neq$ populations distributed in 24 haplogroups.
- Is MSY1 an appropriate marker to study the evolution of these haplogroups? or of the subpopulations in an haplogroup? or of individuals?
- Experiments: Compute all pairwise alignments and construct a phylogenetic tree with BioNJ from our distance matrix.


## Example of a phylogenetic tree



Haplogroup 4 phylogenetic tree.

## Haplogroups tree



Haplogroups tree (2)


## Haplogroups trees comparison



## Conclusion and future works

- A new and specific alignment method for minisatellites maps.
- Extensions: multiple amplification/contraction.

1. arity $>1: \cdots \operatorname{cgg}$ att $\operatorname{tga} \cdots \longrightarrow \cdots \operatorname{cgg}$ att att att $\operatorname{tga} \cdots$
2. order $>1: \cdots \operatorname{cgg} \underline{\text { att tga } \cdots \longrightarrow \cdots \operatorname{cgg} \text { att tga att tga } \cdots \cdot}$

- Applicability for different biological purposes and on different ms.

