Comparison of Minisatellites

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2. Evolutionary model.

3. Arches.

4. Alignment algorithm.

5. Applications.

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Minisatellites (ms): what are they?

- Class of tandem repeat: Satellites, Minisatellites, Microsatellites.
- Example of a tandem repetition:

··· cggcgat cggcgat cggcgat cggcgat cggcgat cggcgat cggcgat ···

- A minisatellite unit measures between 7-100 bp, and their overall length is greater than 0.5kb.
- ms undergo amplifications and contractions \Rightarrow variation in their number of unit.
- ms also undergo point-mutations and homogenization.

Interests in Minisatellites

- Implication in diseases: insulin-dependent diabetes, several cancers, epilepsy, and others [Buard, Jeffreys 97].
- Phylogeny and population studies:
 - Micro-evolution in small range of evolutionary time.
 - Migration of population: out-of-Africa hypothesis [Armour et al 96].
 - Evolution of the Y chromosome [Jobling, Tyler 00].
- Individual or species identification: forensic studies, genetic markers, identification of bacteria [Jeffreys 93]
- Theoretical interest: Why such sequences exist? How do they evolve?

Minisatellites maps

- Difficulty to sequence and assemble minisatellites.
- There is a specific method to obtain the sequence of variants of the unit, **Minisatellite Variant Repeat PCR** [Jeffreys et al. 91]. MVR-PCR yields a **ms map**: a sequence of symbols, each representing a different variant of the repeated unit.
- Example of a minisatellite map:
 - -s = cggcgat cggcgac cggagat cggcgat cggcgat cggagat cgacgat
 - New alphabet: A = cggcgat B = cggcgac C = cggagat D = cgacgat
 - Corresponding map: A B C A A C D

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Ancestor map: AAAAA

Today's variants: A,B,C,D,E

Trace of the evolution in individual 1:

	Α	Α	Α	Α	Α						
Mutation	Α	Α	Α	B	Α						
5 Amplifications	Α	Α	Α	B	B	B	B	B	B	Α	
Mutation	Α	Α	Α	B	B	С	B	B	B	Α	
Mutation	Α	Α	Α	B	B	С	B	D	B	Α	
Amplification	Α	Α	Α	B	B	С	B	D	D	B	Α

Trace of the evolution in individual 2:

	Α	Α	Α	Α	Α	
Mutation	Α	E	Α	Α	Α	
Amplification	Α	E	<u>A</u>	Α	Α	Α
Contraction	Α	E	Α	Α	Α	

We would like to compute an optimal alignment:

Single Step Evolutionary Model

- Unitary operations
 - Amplification(A)/Contraction(C) duplicates/removes a variant which is next to an identical variant.
 - Mutation(M) mutates one variant in another.
 - Insertion(I)/Deletion(D) inserts/removes a variant, but without constraint.
- Non-unitary operations:
 - Generation(G)/Compression(K) generates/compresses an arch from/to its single ancestor variant.
- A cost is associated to each operation. Alignment cost is the sum of its operations costs.

Operations costs

- The model is symmetric: I = D and A = C.
- Observed much higher relative frequency of amplifications and contractions compared to other events \Rightarrow A, C < M, D, I.
- A deletion can also be obtained by a mutation plus a contraction. We have either:
 - 1. Hypothesis 1 (H1): D > M + C and I > A + M
 - 2. Hypothesis 2 (H2): $D \le M + C$ and $I \le A + M$
- **Theorem**: Under H1 or H2, alignment cost is a metric distance.

Problem: Let s and r be two ms maps of length n and m, fi nd an *optimal global* alignment between s and r under the single step evolutionary model.

• Alignment with duplication [Benson 97]:

- Alignment of two sequences that may contain several tandem repeats; operations are substitutions, indels and duplications.

- Differences with our approach: duplication relates the two sequences while for us a duplication is internal to one map; repeated units are unknown.

• Duplication history reconstruction [Benson et Dong 99, Elemento al 01, Tang al 01, Elemento 02, Jaitly 02].

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Let s be a map.

Def: an arch of s is a subsequence of s whose first and last variants are identical.

Def: an arch is single is the internal variants occur only once and complex otherwise.

Ex: s=acacdefda



Def: two arches are **compatible** if they do not cross or share a same first foot or a same last foot.

Generate/Compress an Arch

- Generation and Compression are symmetrical.
- Generate a simple arch: from A to ABCA

First method					Second method						
Amplification Mutation Amplification Mutation Amplification	A A A A A A	A B B B B	B C C	С	Amplification Amplification Amplification Mutation Mutation	A A A A A A	A A A B B	A A A C	A A A		
Mutation	A	B	C	A							

Identical variants at the extremities allow to avoid one mutation.

- Non-commutativity of operations.
- **Theorem:** the optimal generation/compression of an arch uses the largest subset of compatible inner arches.

• Map of size n, arches are subintervals in [1,n].

- Compatibility relation defines an overlap graph.
- **Theorem**: Finding the largest subset of compatible inner arches is equivalent to finding the max stable in an overlap graph.
- Complexity of this procedure is O(V²) with V the number of vertices of the graph [Apostolico et al. 92].
- Maximal number of arches is $\Theta(n^2)$.

Example

acements



How to convert the problem of finding the max subset of compatible arches in the problem of finding the max stable set in an overlap graph

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Dynamic Programming Dependencies



Recurrence

 $\mathscr{A}[i - 1, j - 1] + M(s[i], r[j])$ Mute $\mathscr{A}[i-1,j]+C$ Contract **iff** s[**i** - 1] = s[**i**] **or** s[**i**]=**r**[**j**] $\mathscr{A}[i-1, j] + S$ Delete $\mathscr{A}[i, j-1] + A$ Amplify **iff r**[**j** - 1] = **r**[**j**] **or s**[**i**]=**r**[**j**] $\mathscr{A}[i, j-1] + I$ Insert $\mathscr{A}[1, j] + K(s[1, i])$ Compress an arch $\forall l \in [1, i - 2]$ such that s[l] = s[i] $\mathscr{A}[i, l'] + G(r[l', j])$ Generate an arch \forall l' \in [1, j - 2] such that r[l'] = r[j]

 $\mathscr{A}[i,j] = min$

Theorem: If p = max(n, m), overall complexity is $O(p^4)$ time.

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• MSY1 is a ms on the Y chromosome. Unit length 25bp, 5 variants differ by at most 4 residues. Evolution simplified: no exchange between alleles [Jobling et al 98].

• Maps for 690 men taken in \neq populations distributed in 24 haplogroups.

• Is MSY1 an appropriate marker to study the evolution of these haplogroups? or of the subpopulations in an haplogroup? or of individuals?

• Experiments: Compute all pairwise alignments and construct a phylogenetic tree with BioNJ from our distance matrix.

Example of a phylogenetic tree



Haplogroup 4 phylogenetic tree.



Haplogroups tree (2)



Haplogroups trees comparison

• A new and specific alignment method for minisatellites maps.

- Extensions: multiple amplification/contraction.
 - 1. arity > 1: \cdots cgg <u>att</u> tga $\cdots \longrightarrow \cdots$ cgg **att att att** tga \cdots
 - 2. order > 1: \cdots cgg att tga $\cdots \longrightarrow \cdots$ cgg att tga att tga \cdots

• Applicability for different biological purposes and on different ms.